

RESEARCH OUTPUTS / RÉSULTATS DE RECHERCHE

Comparative Semantics of Feature Diagrams: FFD vs vDFD

Trigaux, Jean-Christophe; Heymans, Patrick; Schobbens, Pierre-Yves; Classen, Andreas

Published in:

Proceedings of the Workshop on Comparative Evaluation in Requirements Engineering (CERE'06) held in conjunction with the 14th IEEE International Requirements Engineering Conference (RE'06)

Publication date:

2006

Document Version

Early version, also known as pre-print

[Link to publication](#)

Citation for published version (HARVARD):

Trigaux, J-C, Heymans, P, Schobbens, P-Y & Classen, A 2006, Comparative Semantics of Feature Diagrams: FFD vs vDFD. in *Proceedings of the Workshop on Comparative Evaluation in Requirements Engineering (CERE'06) held in conjunction with the 14th IEEE International Requirements Engineering Conference (RE'06)*.

General rights

Copyright and moral rights for the publications made accessible in the public portal are retained by the authors and/or other copyright owners and it is a condition of accessing publications that users recognise and abide by the legal requirements associated with these rights.

- Users may download and print one copy of any publication from the public portal for the purpose of private study or research.
- You may not further distribute the material or use it for any profit-making activity or commercial gain
- You may freely distribute the URL identifying the publication in the public portal ?

Take down policy

If you believe that this document breaches copyright please contact us providing details, and we will remove access to the work immediately and investigate your claim.

Comparative semantics of Feature Diagrams: FFD vs. vDFD

Jean-Christophe Trigaux, Patrick Heymans, Pierre-Yves Schobbens, Andreas Classen
University of Namur, Computer Science Department
5000 Namur, Belgium
{jtr,phe,pys,aclassen}@info.fundp.ac.be

Abstract

Feature Diagrams are a popular family of modelling languages used for engineering requirements in software product lines. In our previous research, we advocated the use of formal semantics as an indispensable means to clarify discussions about feature diagrams and to facilitate safe and efficient tool automation. We presented a generic formal semantics for feature diagram languages and criteria to compare them. However, other formal semantics exist. We already informally argued in favour of our semantics which, we think, is more abstract, more concise and not tool dependent. However, some of these claims needed to be further objectified. The purpose of this paper is to compare the semantics proposed by van Deursen and Klint with our own following the methodology of comparative semantics. To be made amenable to comparison, van Deursen and Klint's tool-based definition is first recalled and redefined by correcting some minor mistakes. Their semantics is then mapped to ours through an abstraction function. We then proceed to compare the expressiveness, embeddability and succinctness of both approaches. The study tends to confirm our semantic choices as well as our tool-independent methodology. It also demonstrates that van Deursen and Klint's language is fully expressive and provides various results likely to help tool developers, especially for implementing model transformations.

1 Introduction

Central to the Product Line (PL) paradigm is the modelling and management of *variability*, i.e. *the commonalities and differences in the applications in terms of requirements, architecture, components, and test artifacts* [21]. Variability at the requirement level is commonly modelled through *Feature Diagrams* (FD). In the last decade, research and industry have developed several FD languages. The first and seminal proposal was introduced as part of the FODA method back in 1990 [16]. An example of a FODA FD is

given in Fig. 1. It is inspired from a case study defined in [6] and indicates the allowed combinations of features for a family of systems intended to monitor the engine of a car. As is illustrated, FODA features are nodes of a graph represented by strings and related by various types of edges. On top of the Fig. 1, the node **Monitor Engine System** is called the *root*, or *concept*. The nodes can be mandatory or optional. Optional nodes are represented with a hollow circle above their name, e.g. **Coolant**. In FODA, mandatory nodes are the ones without a hollow circle (in some other syntaxes [12, 23, 22, 7, 10, 8], they are represented with filled circles). The edges are used to progressively decompose nodes into more detailed features. In FODA, there were only *and*- and *xor*-decompositions like illustrated in Fig. 1:

1. *and-decomposition* e.g. between **Monitor Fuel Consumption** and its sons, **Measures** and **Methods**, indicating that they should both be present in all feature combinations where **Monitor Fuel Consumption** is present.
2. *xor-decomposition* where edges are linked by a line segment, as between **Measures** and its sons, **l/km** and **Miles/gallon**, indicating that only one of them should be present in combinations where **Measures** is.

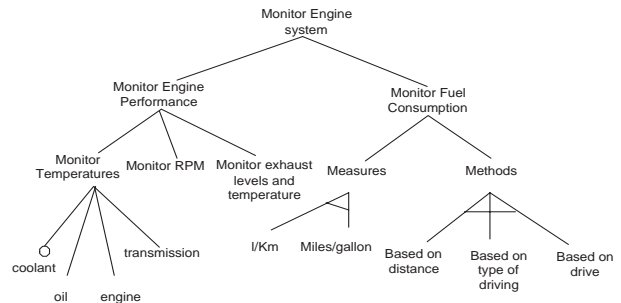


Figure 1. FODA: Monitor Engine System

Since Kang *et al.*'s initial proposal, several extensions have been devised as part of the following methods: FORM [17], FeatureRSEB [12], Generative Programming [10], PLUSS [11], and in the work of the following authors: Riebisch *et al.* [23, 22], van Gurp *et al.* [28], van Deursen and Klint [27], Czarnecki *et al.* [7, 8], Batory [1] and Benavides *et al.* [2].

Most of these authors [17, 12, 10, 23, 22, 28, 11] present their semantics by the way of examples. Still, most of them have argued for an "improved expressiveness". However, without a formal semantics, they have failed to demonstrate it.

In previous publications [4, 26, 25], we have developed and applied a rigorous framework to assess those claims. We have first carried out a comprehensive survey of the informal FD variants. We have generalized their various syntaxes through a generic construction called Free Feature Diagrams (FFD). We gave a formal semantics to FFD, thus providing a (hopefully) unambiguous and very concise definition for all the surveyed FD variants. All formalization choices found a clear answer in the original FODA FD definition, which proved that although informal and scattered throughout many pages, it suffered no ambiguity problem. However, having a proper formal semantics remains extremely important. As remarkably argued in [13], formal semantics is the best way to avoid ambiguities and to start building safe automated reasoning tools. Without a formal semantics, new FD languages might continue to proliferate on the basis of shallow or erroneous motivations, leading to interpretation and interoperability problems.

In [4, 26, 25], we argued that FFD contribute to improve the definition, understanding, comparison and reliable implementation of FD languages. In particular, we have highlighted some subtle points in the interpretation of FD. Additionally, we have defined the main decision problems that a FD tool should implement, i.e. we gave a specification of such tools, and subsequently, we studied their minimal execution time.

Some authors have also started to better define their semantics [27, 8, 7, 1]. However, we found that these semantics are less general, abstract and concise. These approaches typically transform FD to other formalisms (for which tools exist). This naturally gives a more complex transformation and a less abstract semantics. It has the dubious advantage that the transformation is correct by definition. On the contrary, we believe that tools should be built or chosen according to a natural, carefully chosen and well-studied semantics. Our approach is justified by our goals: make fundamental semantic issues of FD languages explicit in order to study their properties and rigorously evaluate them before adopting them or implement CASE tools.

We are now in position to proceed to our next phase of study sketched in [26, 25]: given well-defined FD lan-

guages, we can start a meaningful discussion of their merits, following the well-established scientific method called *comparative semantics*. In this paper, we compare the semantics of FFD with the one defined by van Deursen and Klint in [27] which is apparently the first formal semantics of FD to have been published. For brevity, we call their variant of FD, vDFD (van Deursen and Klint's Feature Diagrams¹).

This paper is structured as follows. In Sec. 2, we will describe the method that we use to compare formal semantics. FFD are then briefly presented in Sec. 3. In Sec. 4, we will recall the semantics of vDFD given in [27] and compare it with our own. The comparison of both formalisms appears in Sec. 5. Finally, related works are examined in Sec. 6 and conclusions are given in Sec. 7.

2 Research method

A proper definition of a formal semantics is preferably done in several steps. The first step in this chain is to have a bidirectional conversion from the *concrete syntax* (what the user sees), to abstract data structures, called *abstract syntax*. Indeed, the concrete syntax is usually too cumbersome to be reasoned on efficiently, be it by humans or tools.

The abstract syntax is usually a graph. It is thus essential to specify exactly what are the allowed graphs. There are two common ways to provide this information: (1) *mathematical notation* (set theory) or (2) *meta-model* (usually, a UML Class Diagram complemented with OCL constraints). In this paper, we follow [13] and adopt the former for its improved conciseness, rigour and suitability for performing mathematical proofs. The latter might be more readable and facilitate some implementation tasks such as building a repository.

The second step is to provide a formal semantics, i.e. a function from the graphs above to a mathematical structure chosen for being as close as possible to our intuitive understanding. This structure forms the *semantic domain*. In [3] (and in Def. 3.5 of this paper) we proposed as semantic domain the one of *Product Lines (PL)* defined as set of *products*, where a product is characterized by the primitive features it includes.

The works to which we compare often do not follow this methodology, but are amenable to it. For instance, [1] defines a transformation to grammars and propositional formulae. Fortunately, these two formalisms are provided with a standard semantics, so that we can obtain a semantics by composing the transformation followed by the standard semantics. In all similar approaches, the semantic domain of the formalisms used were not designed for features. They are thus usually less abstract: they keep too much syntactic

¹van Deursen and Klint's Feature Diagrams are restricted graphs in the sense that they are trees, except for their leaves which can be shared.

information, so that fewer diagrams are considered equivalent. To discard this syntactic information, we must introduce *abstraction functions* between semantic domains (see Fig. 2). As we progress in our comparative semantics study, we will construct a category of semantic domains linked by abstraction functions. The comparative semantics of specification languages [19] of logic programming, of concurrent programming [9] or of coordination languages [5], for instance, is already well developed and integrates developments and tools from different languages.

Given this, we can then evaluate which FD languages are more expressive, succinct, or natural [3] more rigorously. The technical tool for this study are *translations* or *embeddings* (see Fig. 2), i.e. transformations (functions between abstract syntaxes) that preserve semantics.

As illustrated in Fig. 2, for each FD language, say X , with an already existing formal semantics, their semantic domain (XSD) and abstract syntax (XFD) should be compared with our semantic domain (PL, see Def. 3.5) and the abstract syntax of FFD (Def. 3.1), respectively, in order to derive abstraction functions and embeddings. Moreover, these translations compose, so that it is useful to have our category of semantic domains, and look at its shape so that most results follow by composition.

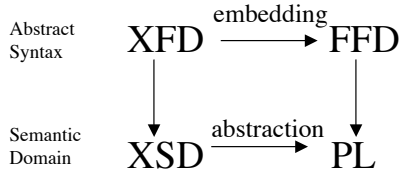


Figure 2. Comparative Semantics Approach

3 Free Feature Diagrams: FFD

In this section, for the paper to be self-contained, we recall *Free Feature Diagrams* (FFD) from [26, 25], a parametric construction that generalizes the syntax of FD languages and represents a family of FD languages. We define it according to our research method.

3.1 Abstract syntax

All FD are graphs whose nodes are features. They were drawn as (boxed) strings in the concrete FD languages. Many authors use the word “feature” ambiguously, sometimes to mean a node, sometimes a leaf (a node that is not further decomposed), and sometimes as a real-world feature. Here, we use the term “feature” as a synonym of “node”.

We further distinguish “primitive” and “compound features”²³. *Primitive features* are “features” that are of interest *per se*, and that will influence the final product. On the contrary, *compound features* are just intermediate nodes used for decomposition. For generality, we leave it to the modeler to define which nodes in the FD have such a purpose. Primitive features are thus not necessarily equivalent to leaves, though it is the most common case.

Decomposition edges relate a *father* node f to a *son* node s and are noted $f \rightarrow s$.

FD languages vary in several respects: (1) do they consider FD as trees or DAG, (2) what are the allowed types of operators on nodes, (3) what are the allowed types of graphical constraints, (4) what are the allowed textual constraints. These are the parameters (GT,NT,GCT,TCL) of FFD:

- GT (Graph Type) is either DAG or TREE.
- NT (Node Type) is a set of Boolean functions (operators), at most one per arity. E.g.: *and* is the set of operators and_s , one for each arity s , that return true iff all their s arguments are true. Similarly, *or* (resp. *xor*) is the set of operators or_s (resp. xor_s) that return true iff some (resp. exactly one) of their s arguments is true. Operators opt_s in *opt* always return true. Operators $vp_s(i..j)$ in *card* return true iff at least i and at most j of their arguments are true. NT is usually some combination of those sets.
- GCT (Graphical Constraint Type) is a binary Boolean operator. E.g.: *Requires* (\Rightarrow) or *Mutex* (\cap).
- TCL (Textual Constraint Language) is a subset of the language of Boolean formulae where the predicates are the nodes of the FD. The sublanguage used in FODA FD, “Composition rules” [16, p.71] is: $CR ::= p_1(\text{requires} \mid \text{mutex})p_2$ where p_1, p_2 are primitive features.

The syntactic domain of a particular FD language can be defined simply by providing values for these parameters. For example, the abstract syntax of FODA FD is defined as $FFD(\text{TREE}, and \cup xor \cup \{opt_1\}, \emptyset, CR)$. In [4], the abstract syntax of other FD variants, including [17, 12, 10, 23, 22, 28, 11] is defined similarly. As we will see in Sec. 4.2, van Deursen and Klint’s abstract syntax is defined as $FFD(\text{DAG}, and \cup xor \cup or \cup \{opt_1\}, \emptyset, CR')$ where $CR' ::= p_1(\text{requires} \mid \text{excludes})p_2 \mid (\text{include} \mid \text{exclude})p$.

The semantics is defined only once for FFD [26, 25], reproduced in Sec. 3.2. The formal semantics of a particular FD language defined through FFD thus comes for free.

²We adopt the terminology of [1].

³“Compound features” are also called “decomposable features”.

Definition 3.1 (Free Feature Diagram, or FFD) A FFD $d \in FFD(GT, NT, GCT, TCL) = (N, P, r, \lambda, DE, CE, \Phi)$ where:

- N is its set of nodes;
- $P \subseteq N$ is its set of primitive nodes;
- $r \in N$ is the root of the FD, also called the concept;
- $\lambda : N \rightarrow NT$ labels each node with an operator from NT ;
- $DE \subseteq N \times N$ is the set of decomposition edges; $(n, n') \in DE$ will rather be noted $n \rightarrow n'$;
- $CE \subseteq N \times GCT \times N$ is the set of constraint edges;
- $\Phi \subseteq TCL$ are the textual constraints.

FFD collect whatever can be drawn. So, FD have additional minimal well-formedness conditions.

Definition 3.2 (Feature Diagram, or FD) A FD is a FFD where:

1. Only r has no parent: $\forall n \in N. (\nexists n' \in N. n' \rightarrow n) \Leftrightarrow n = r$.
2. DE is acyclic: $\nexists n_1, \dots, n_k \in N. n_1 \rightarrow \dots \rightarrow n_k \rightarrow n_1$.
3. If $GT = TREE$, DE is a tree: $\nexists n_1, n_2, n_3 \in N. n_1 \rightarrow n_2 \wedge n_3 \rightarrow n_2 \wedge n_1 \neq n_3$.
4. Nodes are labelled with operators of the appropriate arity: $\forall n \in N. \lambda(n) = op_k \wedge k = \#\{(n, n') | n \rightarrow n'\}$.

3.2 Semantics

A formal semantics is given by a function from the syntactic domain of a language to a semantic domain [13]. The syntactic domain was given in Def. 3.1 and 3.2. Here, after some preliminary definitions, we define the semantic domain as the set of product lines (PL) (Def. 3.5, point 3) and then the semantic function (Def. 3.5, point 4).

The notion of model for a FD was introduced in [16, p.64], with the examples of models of X10 terminals.

Definition 3.3 (Model) A model of a FD is a subset of its nodes: $M = \mathcal{P}N$.

Definition 3.4 (Valid model) [16, p.70] A model $m \in M$ is valid for a $d \in FD$, noted $m \models d$, iff:

1. The concept is in: $r \in m$
2. The meaning of nodes is satisfied: If a node $n \in m$, and n has sons s_1, \dots, s_k and $\lambda(n) = op_k$, then $op_k(s_1 \in m, \dots, s_k \in m)$ must evaluate to true.

3. The model must satisfy all textual constraints: $\forall \phi \in \Phi. m \models \phi$, where $m \models \phi$ means that we replace each node name n in ϕ by the truth value of $n \in m$, evaluate ϕ and get true. For example, if we call ϕ_1 the CR constraint p_1 requires p_2 , we say that $m \models \phi_1$ when $p_1 \in m \Rightarrow p_2 \in m$ is true.
4. The model must satisfy all graphical constraints: $\forall (n_1, op_2, n_2,) \in CE, op_2(n_1 \in m, n_2 \in m)$ must be true.
5. If s is in the model and s is not the root, one of its parents n , called its justification, must be too: $\forall s \in N. s \in m \wedge s \neq r : \exists n \in N : n \in m \wedge n \rightarrow s$.

Definition 3.5 (Product and Product Line, or PL) We define:

1. A product c is a set of primitive nodes: $c \in \mathcal{P}P$.
2. The product named by a model m , noted $\llbracket m \rrbracket$, is $m \cap P$.
3. A product line (PL) pl is a set of products: $pl \in \mathcal{P}\mathcal{P}P$.
4. The product line of a FD d consists of the products named by its valid models: $\llbracket d \rrbracket = \{m \cap P | m \models d\}$.

4 van Deursen and Klint's Feature Diagrams: vDFD

van Deursen and Klint have formalized FD, by providing [27]:

1. A syntax presented as a feature description language that we will call vDFD. The authors also provided a feature definition (Def. 4.1) and a grammar for vDFD (Def. 4.3);
2. A semantics presented as a feature diagram algebra defining various sets of rules manipulating vDFD:
 - (a) Normalization rules (\mathcal{N}) to eliminate duplicate features and degenerate cases of the various constructs (Def. 4.5);
 - (b) Variability rules to count the number of products allowed in a FD;
 - (c) Expansion rules (\mathcal{E}) to expand a normalized feature expression into a disjunctive normal form (Def. 4.6);
 - (d) Satisfaction rules (\mathcal{S}) to determine which feature expressions in disjunctive normal form satisfy the feature constraints (Def. 4.7);

In Fig. 3, we show the sequence of these transformations ($\mathcal{N}, \mathcal{E}, \mathcal{S}$) on vDFD as proposed in [27]. An alternative sequence of transformations ($\mathcal{N}', \mathcal{E}', \mathcal{S}'$) on vDFD is shown which refers to the small corrections we propose for each of them. In the next sections, we will give the reader more details on this. First, we will present the formalization for vDFD proposed in [27] (Sec. 4.1), then we will discuss their abstract syntax (Sec. 4.2) and revisit their semantics (Sec. 4.3). Finally, we compare the revisited semantics with our own (Sec. 5).

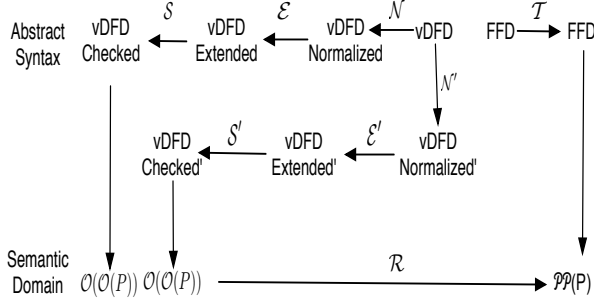


Figure 3. van Deursen and Klint's semantics

4.1 van Deursen and Klint's original definition

The primary objective of van Deursen and Klint was to reason on FD using a textual representation rather than a graphical one. Further requirements for this representation were [27]:

1. to contain all the information contained in the graphical form,
2. to be suited for automatic reasoning.

To satisfy these requirements, the authors first produced a feature definition (Def. 4.1). Then they proposed a grammar for their textual FD (Def. 4.3). Finally, they proposed rules to reason on this representation. Rules are used to check the consistency of the representation. \mathcal{N} and \mathcal{E} are used to generate a normal form (syntactic consistency). \mathcal{S} are used to check constraints satisfaction (semantic consistency). Rules for computing variability are also defined but are not relevant here as they do not influence the semantics. All these rules are defined and justified in [27]. To present these definitions we reuse the variable naming convention proposed by van Deursen and Klint (Table 1). All the reasoning proposed by the authors is based on their disjunctive normal form (Def. 4.4).

Definition 4.1 (Feature definition) A *feature definition* [27, p.4] is a feature name followed by “:” and a feature expression (Def. 4.2)

Definition 4.2 (Feature expression) A *feature expression* [27, p.4] can consist of

1. an atomic feature,
2. a composite feature: a named feature whose definition appears elsewhere,
3. an optional feature: a feature expression followed by “?”,
4. mandatory features: a list of feature expressions enclosed in `all()`,
5. alternative features: a list of feature expressions enclosed in `one-of ()`,
6. non-exclusive selection of features: a list of feature expressions enclosed in `more-of ()`,
7. a default feature value: `default =` followed by an atomic feature,
8. features of the form `...,` indicating that a given set is not completely specified.

Definition 4.3 (vDFD Grammar) A vDFD Grammar [27, p.6] is defined by:

$[A - Z][a - zA - Z0 - 9]^*$	\rightarrow FeatureName
$[a - z][a - zA - Z0 - 9]^*$	\rightarrow AtomicFeature
FeatureDefinition*	\rightarrow FeatureDiagram
Constraint*	\rightarrow FeatureDiagram
FeatureName:	\rightarrow FeatureDefinition
FeatureExpr	\rightarrow FeatureList
{ FeatureExpr , }+	\rightarrow FeatureList
all(FeatureList)	\rightarrow FeatureExpr
one-of (FeatureList)	\rightarrow FeatureExpr
more-of (FeatureList)	\rightarrow FeatureExpr
FeatureName	\rightarrow FeatureExpr
AtomicFeature	\rightarrow FeatureExpr
FeatureExpression ?	\rightarrow FeatureExpr
default = AtomicFeature	\rightarrow FeatureExpr
...	\rightarrow AtomicFeature
DiagramConstraint	\rightarrow Constraint
UserConstraint	\rightarrow Constraint
AtomicFeature requires	\rightarrow DiagramConstraint
AtomicFeature excludes	\rightarrow DiagramConstraint
AtomicFeature	\rightarrow DiagramConstraint
include AtomicFeature	\rightarrow UserConstraint
exclude AtomicFeature	\rightarrow UserConstraint

Definition 4.4 (Disjunctive normal form) A *disjunctive normal form* [27, p.9] is a `one-of` feature expression with only `all` feature expressions as arguments themselves with only atomic features as arguments. A disjunctive normal form is an expression of the form: `one-of(all(A_{11}, \dots, A_{1n_1}), ..., all(A_{m1}, \dots, A_{mn_m}))`

Meta-Variable	Type
F	FeatureExpr
Fs,Fs',Fs''	{FeatureExpr “,”}*
Ft	{FeatureExpr “,”}+
A,A1,A2	AtomicFeature
C	Constraint
Cs,Cs'	Constraint*

Table 1. Variable naming conventions
(adapted from [27, p.7])

This disjunctive normal form indicates explicitly all possible feature combinations. It is obtained by applying the normalization (\mathcal{N}) and expansion rules (\mathcal{E}). For instance, (\mathcal{N}_2) removes duplicate features; (\mathcal{N}_8) transforms a `one-of` expression containing one optional feature into an optional `one-of` expression; (\mathcal{E}_4) translates an `all()` expression containing a `more-of` expression into three cases: one with the first alternative, one with the first alternative and the remaining `more-of` expression, and one with only the remaining `more-of` expression.

Definition 4.5 (Normalization rules) *The set of normalization rules [27, p.7] is $\mathcal{N} = \{\mathcal{N}_1, \dots, \mathcal{N}_{12}\}$:*

- (\mathcal{N}_1) $Fs, F, Fs', F?, Fs'' = Fs, F, Fs', Fs'$
- (\mathcal{N}_2) $Fs, F, Fs', F, Fs'' = Fs, F, Fs', Fs''$
- (\mathcal{N}_3) $F?? = F?$
- (\mathcal{N}_4) $\text{all}(F) = F$
- (\mathcal{N}_5) $\text{all}(Fs, \text{all}(Ft), Fs') = \text{all}(Fs, Ft, Fs')$
- (\mathcal{N}_6) $\text{one-of}(F) = F$
- (\mathcal{N}_7) $\text{one-of}(Fs, \text{one-of}(Ft), Fs') = \text{one-of}(Fs, Ft, Fs')$
- (\mathcal{N}_8) $\text{one-of}(Fs, F?, Fs') = \text{one-of}(Fs, F, Fs')?$
- (\mathcal{N}_9) $\text{more-of}(F) = F$
- (\mathcal{N}_{10}) $\text{more-of}(Fs, \text{more-of}(Ft), Fs') = \text{more-of}(Fs, Ft, Fs')$
- (\mathcal{N}_{11}) $\text{more-of}(Fs, F?, Fs') = \text{more-of}(Fs, F, Fs')?$
- (\mathcal{N}_{12}) $\text{default} = A = A$

Definition 4.6 (Expansion rules) *The set of expansion rules [27, p.9] is $\mathcal{E} = \{\mathcal{E}_1, \dots, \mathcal{E}_4\}$:*

- (\mathcal{E}_1) $\text{all}(Fs, F?, Ft) = \text{one-of}(\text{all}(Fs, F, Ft), \text{all}(Fs, Ft))$
- (\mathcal{E}_2) $\text{all}(Ft, F?, Fs) = \text{one-of}(\text{all}(Ft, F, Fs), \text{all}(Ft, Fs))$
- (\mathcal{E}_3) $\text{all}(Fs, \text{one-of}(F, Ft), Fs') = \text{one-of}(\text{all}(Fs, F, Fs'), \text{all}(Fs, \text{one-of}(Ft), Fs'))$
- (\mathcal{E}_4) $\text{all}(Fs, \text{more-of}(F, Ft), Fs') = \text{one-of}(\text{all}(Fs, F, Fs'), \text{all}(Fs, F, \text{more-of}(Ft), Fs'), \text{all}(Fs, \text{more-of}(Ft), Fs'))$

On this disjunctive normal form, satisfaction rules (Def. 4.7) are applied to eliminate products that do not satisfy the constraints. These satisfaction rules use the two following functions (not explicitly defined in [27]):

- $\text{isElement} : \text{AtomicFeature} \times \text{FeatureExpression} \rightarrow \mathcal{B}$ which determines whether the *AtomicFeature* is contained in the *FeatureExpression* or not.
- $\text{sat} : \text{FeatureExpr} \times \text{Constraints} \rightarrow \mathcal{B}$ which determines whether the *FeatureExpr* satisfies the *Constraints* or not.

If no constraint is applicable to the feature expression then (\mathcal{S}_9) is used. Otherwise, binary constraints (`requires`, `excludes`) and unary constraints (`include`, `exclude`) are respectively handled by :

- (\mathcal{S}_1) and (\mathcal{S}_2) for `excludes`;
- (\mathcal{S}_3) and (\mathcal{S}_4) for `requires`;
- (\mathcal{S}_5) and (\mathcal{S}_6) for `include`;
- (\mathcal{S}_7) and (\mathcal{S}_8) for `exclude`.

For instance, (\mathcal{S}_6) defines that the *sat* function must return *false* if the constraint *Include A* is applicable and if the atomic feature *A* is not an element of the *FeatureExpr Ft*.

Definition 4.7 (Satisfaction rules) *The set of satisfaction rules [27, p.13] is $\mathcal{S} = \{\mathcal{S}_1, \dots, \mathcal{S}_9\}$, where $|$ means OR:*

- (\mathcal{S}_1) $\frac{\text{isElement}(A2, Fs) \mid \text{isElement}(A2, Fs') = \text{true}}{\text{sat}(\text{all}(Fs, A1, Fs'), Cs \text{ A1 excludes A2 Cs'}) = \text{false}}$
- (\mathcal{S}_2) $\frac{\text{isElement}(A2, Fs) \mid \text{isElement}(A2, Fs') = \text{false}}{\text{sat}(\text{all}(Fs, A1, Fs'), Cs \text{ A1 excludes A2 Cs'}) = \text{sat}(\text{all}(Fs, A1, Fs'), Cs Cs')}$
- (\mathcal{S}_3) $\frac{\text{isElement}(A2, Fs) \mid \text{isElement}(A2, Fs') = \text{false}}{\text{sat}(\text{all}(Fs, A1, Fs'), Cs \text{ A1 requires A2 Cs'}) = \text{false}}$
- (\mathcal{S}_4) $\frac{\text{isElement}(A2, Fs) \mid \text{isElement}(A2, Fs') = \text{true}}{\text{sat}(\text{all}(Fs, A1, Fs'), Cs \text{ A1 requires A2 Cs'}) = \text{sat}(\text{all}(Fs, A1, Fs'), Cs Cs')}$
- (\mathcal{S}_5) $\frac{\text{isElement}(A, Ft) = \text{true}}{\text{sat}(\text{all}(Ft), Cs \text{ include A Cs'}) = \text{sat}(\text{all}(Ft), Cs Cs')}$

$$\begin{aligned}
(\mathcal{S}_6) \quad & \frac{isElement(A, Ft) = false}{sat(all(Ft), Cs \text{ include } A \text{ } Cs') = false} \\
(\mathcal{S}_7) \quad & \frac{isElement(A, Ft) = true}{sat(all(Ft), Cs \text{ exclude } A \text{ } Cs') = false} \\
(\mathcal{S}_8) \quad & \frac{isElement(A, Ft) = false}{sat(all(Ft), Cs \text{ exclude } A \text{ } Cs') = sat(all(Ft), Cs \text{ } Cs')} \\
(\mathcal{S}_9) \quad & sat(all(Ft), Cs) = true
\end{aligned}$$

4.2 Abstract syntax

van Deursen and Klint have clearly defined the concrete syntax of the vDFD and the various operations to manipulate the allowed expressions. In terms of FFD, we understand that their abstract syntax is $FFD(DAG, and \cup xor \cup or \cup \{opt_1\}, \emptyset, CR')$ where $CR' ::= p_1(\text{requires} \mid \text{excludes})p_2 \mid (\text{include} \mid \text{exclude})p$. The mapping between their concrete syntax and our abstract syntax is presented in Table 2. vDFD is a restricted graph as the different operators construct a tree, except for the leaves which can be shared by several fathers (Def. 4.8).

Concrete Syntax	Abstract Syntax
all	$\rightarrow and$
one-of	$\rightarrow xor$
more-of	$\rightarrow or$
?	$\rightarrow opt_1$

Table 2. vDFD to FFD

Definition 4.8 (van Deursen and Klint's FD) A vDFD is a FFD $= (N, P, r, \lambda, DE, CE, \Phi)$ where:

$$\forall n1, n2, n3 \in N. n1 \rightarrow n3 \wedge n2 \rightarrow n3 \wedge n1 \neq n2 \implies \nexists n4 \in N. n3 \rightarrow n4$$

4.3 Semantics

van Deursen and Klint have defined rewriting rules that lead to disjunctive normal form (Def. 4.4). From this normal form, the semantics is trivial and corresponds to an ordered list of ordered lists of primitive features, that we note $O(O(P))$.

Definition 4.9 (vDFD Semantics) The semantics of a vDFD (vdfg) is a function $\mathcal{L} : vDFD \rightarrow O(O(P))$ where $\mathcal{L}(vdfg) = \mathcal{S}(\mathcal{E}(\mathcal{N}(vdfg)))$.

Nevertheless, we have discovered undesirable semantics due to missing rules. Consequently, we provide some additional rules and justify them:

- The rule in \mathcal{N}_1 is not sufficient to avoid feature lists that combine mandatory and optional features. Indeed, a feature list such as $Fs, F?, Fs', F, Fs''$ (where $F?$ and F are switched wrt the rule \mathcal{N}_1) would be considered normalized. The set of normalisation rules should be corrected by adding one simple rule (Def. 4.10);
- The set of expansion rules is not sufficient to produce a correct disjunctive normal form. Indeed, terms of the form a or $\text{one-of}(Fs)$ or $\text{all}(Fs)$ are allowed. In order to respect the intentions of the authors we extend \mathcal{E} (Def. 4.11);
- The satisfaction function (sat) is never explicitly called. Consequently, we propose one rule (Def. 4.12) that calls this function and eliminates invalid products (products which do not satisfy the constraints).

Definition 4.10 (Normalization rules) The normalization rules are a set of rules $\mathcal{N}' = \mathcal{N} \cup \{\mathcal{N}_{13}\}$ where

$$(\mathcal{N}_{13}) \quad Fs, F?, Fs', F, Fs'' = Fs, F, Fs', Fs''$$

Definition 4.11 (Expansion rules) The expansion rules are a set of rules $\mathcal{E}' = \mathcal{E} \cup \{\mathcal{E}_5, \mathcal{E}_6\}$ where

$$\begin{aligned}
(\mathcal{E}_5) \quad A &= all(A) \\
(\mathcal{E}_6) \quad all(Fs) &= one-of(all(Fs))
\end{aligned}$$

Definition 4.12 (Satisfaction call rule) The satisfaction rules are a set of rules $\mathcal{S}' = \mathcal{S} \cup \{\mathcal{S}_{10}\}$ where

$$(\mathcal{S}_{10}) \quad \frac{sat(all(Fs), Cs)}{sat(one-of(Fs', all(Fs), Fs''), Cs) = sat(one-of(Fs', Fs''), Cs)}$$

Having revisited van Deursen and Klint's rewriting rules we need to redefine the semantics function (Def. 4.13).

Definition 4.13 (Revisited vDFD Semantics) The revisited semantics function of van Deursen and Klint is defined as: $\mathcal{L}' : vDFD \rightarrow O(O(P))$ where $\mathcal{L}'(vdfg) = \mathcal{S}'(\mathcal{E}'(\mathcal{N}'(vdfg)))$

5 Comparison

We will now compare the redefined semantics of vDFD with the one of FFD (Section 5.1) and subsequently define further comparison criteria and examine the languages with respect to those criteria (Section 5.2).

5.1 Comparative Semantics

Even with the modifications of vDFD that we proposed, when we compare their semantics (Def. 4.13) with the one of FFD (Sec. 3.2) two points of divergence appear:

- First, the semantic domains are different (see Fig. 3):
 1. On the one hand, the semantics of vDFD translates a vDFD expression into another expression in disjunctive normal form which is an ordered list of ordered lists of atomic features $O(O(P))$.
 2. On the other hand, the semantic domain of FFD is a set of sets of primitive features $\mathcal{P}(\mathcal{P}(P))$.
- Second, van Deursen and Klint’s semantics always gives preference to inclusion in terms (not in constraints) and thus behaves like a semantics based on edges rather than on nodes.

Let us examine these two points in turn. Although the semantic domains are different, they can be easily related by an abstraction function. Indeed, in [27], atomic features directly correspond to primitive features. Hence, the remaining difference is the one between the mathematical structures *list* and *set*. The order of atomic features is important in van Deursen and Klint’s semantics. For instance, the textual vDFD expression `one-of(all(A, B), all(B, A))` contains two products: `all(A, B)` and `all(B, A)`. In FFD’s semantics, only one product would be part of PL: $\{A, B\}$. Consequently, we can define an obvious abstraction function R (definition 5.1) that simply abstracts away order from van Deursen and Klint’s semantic domain and directly maps it to our semantic domain. We do not know if this notion of order between features was deliberate or not, but intuitively we consider that two products with the same features should be identical. However, ordering features could be relevant, for example, when each feature corresponds to one transformation (on code or models) [24] and these transformations do not produce the same result if they are applied in a different order.

Definition 5.1 (Semantic Domain Abstraction \mathcal{R})

$\mathcal{R} : O(O(P)) \rightarrow \mathcal{P}(\mathcal{P}(P))$

$$\mathcal{R}(\text{all}(f_n, \dots, f_m), \dots, \text{all}(f_{n'}, \dots, f_{m'})) = \{\{f_n, \dots, f_m\}, \dots, \{f_{n'}, \dots, f_{m'}\}\}$$

Nevertheless, this abstraction function is not sufficient to find an equivalence between our semantics and van Deursen and Klint’s semantics. Indeed, the latter implicitly always gives preference to inclusion in terms, which is also a characteristic of edge-based semantics (as opposed to node-based semantics). For instance, contrary to a node-based semantics and classical intuition, an edge-based semantics

for the FD illustrated in Fig. 4 will consider as valid the products $\{A, B, C, D, E\}$ and $\{A, B, C, E, F\}$. The solution to find a semantic equivalence between both semantics (Theorem 5.1) is to apply a preliminary transformation \mathcal{T} on the FFD representation of a textual vDFD. The idea behind the transformation \mathcal{T} is to replace each atomic feature shared by several fathers with one *and*-node for each incoming edge, each of these *and*-nodes having only one son which is the corresponding atomic feature. In our concrete notation, to obtain an edge-based semantics, we add one mandatory circle (*and*-node with one son) for each incoming edge of a shared feature (see Fig. 4).

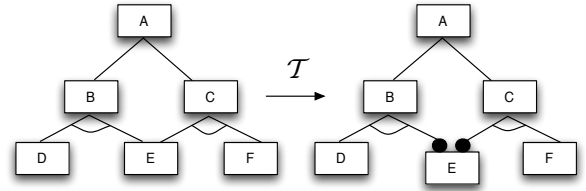


Figure 4. From node-based to edge-based Semantics

Theorem 5.1

$$\forall t : t \in \text{vDFD} : \llbracket \mathcal{T}(t) \rrbracket = \mathcal{R}(\mathcal{L}'(t))$$

5.2 Comparison Criteria

Now that we have aligned the definitions of the two languages, we can start to compare them based on various formal criteria. Three important criteria that we have already studied for other notations [26] are expressiveness, embeddability and succinctness.

FFD and vDFD can not be compared directly with these comparison criteria. Indeed, every possible FD language defined through FFD potentially evaluates differently. Ideally, every FFD-based language should be compared with vDFD separately, which is by far out of the scope of this paper. Hence, in the following, we will introduce the criteria and, for each of them, discuss the comparison between vDFD and one or more representative members from the FD language family that we already studied. More studies could come in the future to complement these results.

5.2.1 Expressiveness

We use the formal, well established, definition of *expressiveness* of a language, as the part of its semantic domain that it can express.

Definition 5.2 (Expressiveness) The expressiveness of a language \mathcal{L} is the set $E(\mathcal{L}) = \{\llbracket D \rrbracket \mid D \in \mathcal{L}\}$, also noted $\llbracket \mathcal{L} \rrbracket$. A language \mathcal{L}_1 is more expressive than a language \mathcal{L}_2 if $E(\mathcal{L}_1) \supset E(\mathcal{L}_2)$. A language \mathcal{L} with semantic domain \mathbb{M} (i.e. its semantics is $\llbracket \cdot \rrbracket : \mathcal{L} \rightarrow \mathbb{M}$) is expressively complete if $E(\mathcal{L}) = \mathbb{M}$.

The usual way to prove that a language \mathcal{L}_2 is at least as expressive as \mathcal{L}_1 is to provide a translation from \mathcal{L}_1 to \mathcal{L}_2 :

Definition 5.3 (Translation) A translation is a total function $T : \mathcal{L}_1 \rightarrow \mathcal{L}_2$ that is correct, i.e. preserves semantics: $\llbracket T(D_1) \rrbracket = \llbracket D_1 \rrbracket$.

Given that we have the domain abstraction function \mathcal{R} , we can consider that the semantic domain of vDFD are also product lines (Def. 3.5). Thus, every vDFD expresses a PL. Now, we can ask the converse question: can every PL be expressed by a vDFD? Stated otherwise: are vDFD fully expressive?

In [26], we have examined this question for several FD languages [16, 17, 12, 10, 11, 23, 22, 28]. The definitions of those languages are recalled in Table 3⁴.

Name	GT	NT	GCT	TCL
OFT	TREE	$and \cup xor \cup \{opt_1\}$	\emptyset	CR
OFD	DAG	$and \cup xor \cup \{opt_1\}$	\emptyset	CR
RFD	DAG	$and \cup xor \cup or \cup \{opt_1\}$	$\{\Rightarrow, \}$	CR
EFD	DAG	$card \cup \{opt_1\}$	$\{\Rightarrow, \}$	CR
GPFT	TREE	$and \cup xor \cup or \cup \{opt_1\}$	\emptyset	CR
PFT	TREE	$and \cup xor \cup or \cup \{opt_1\}$	$\{\Rightarrow, \}$	\emptyset
vDFD	DAG	$and \cup xor \cup or \cup \{opt_1\}$	\emptyset	CR'

Table 3. FD variants defined on top of FFD

If we ignore the graphical and textual constraints, that is, the last two parameters in FFD, we can prove formally [26] that tree languages (OFT [16], GPFT [10], PFT [11]) are not fully expressive. However, DAG languages (OFD [17], EFD [23, 22], RFD [12]) are fully expressive. More precisely, our results show that the disjunction of features cannot be expressed in OFT. In [12], Griss *et al.* have proposed to solve this problem by (1) adding *or*-nodes, and (2) considering FD as single-rooted DAG rather than trees. In [26], we proved that the second extension alone guarantees full expressiveness while adding *or*-nodes only does not.

In vDFD, we have *or*-nodes but we do not have DAGs, or at least just a restricted form of DAG where only the sharing of leaf nodes is allowed. Studying the expressiveness of vDFD thus requires specific treatment.

The operators that manipulate the vDFD expressions must always have at least one operand. Therefore, vDFD

⁴For simplicity, we have given abbreviated names to those languages which are most often different from their original names, if any. References are given in the text

expressions are expressively *incomplete* with respect to PL as the *empty PL* (i.e. the PL containing no product i.e. $\{\}$) and the *base PL* (i.e. the PL containing one product in which no feature has been selected i.e. $\{\{\}\}$) can not be expressed in their disjunctive normal form. If we add the vDFD textual constraints, these two products can be expressed since preference is given to the constraints. An empty PL can be expressed by: a normal form one-of ($all(A)$) and a constraint “exclude A ”, where A is an atomic feature. A base PL can be expressed by: a normal form one-of ($all(A?)$) and a constraint “exclude A ” where A is an atomic feature.

The consequence of this result is that we now know that vDFD equipped with constraints (at least exclude constraints) can be used to express any PL. This is interesting because vDFD are supported by a tool environment and so in theory all FD languages with PL semantics can also be supported by this environment, provided that forward and backward translations between vDFD and the other languages are implemented. We now discuss the practical feasibility of these translations with the two remaining criteria.

5.2.2 Embeddability

In [26], we have proposed a definition of graphical embeddability (Def. 5.4) which generalizes the definition of embeddability for context-free languages (here, simplified):

Definition 5.4 (Graphical embeddability) A graphical language \mathcal{L}_1 is embeddable into \mathcal{L}_2 iff there is a translation $T : \mathcal{L}_1 \rightarrow \mathcal{L}_2$ that is node-controlled [15] : T is expressed as a set of rules of the form $D_1 \rightarrow D_2$, where D_1 is a diagram containing a defined node or edge n , and all possible connections with this node or edge. Its translation D_2 is a subgraph in \mathcal{L}_2 , plus how the existing relations should be connected to nodes of this new subgraph.

Given this definition, we only need to look at the (graph-based) abstract syntax of FD languages to study their embeddability.

All tree FD languages are clearly embeddable into vDFD. For DAG languages, this is not the case because we can have sharing of intermediate nodes in the graph. However, if we consider sublanguages that restrict the sharing to leaves, we just have to apply the linear transformation \mathcal{T} (to guarantee edge-based semantics) and from Theorem 5.1 we can directly infer that we have an embedding. Finally, vDFD are embeddable into RFD.

Embeddings are of practical relevance because they ensure that there exists a transformation from one language to the other which preserves the whole shape of the models. This way, traceability between the two models is greatly facilitated and tool interoperability is made more transparent.

Embedding results must however be completed by examining the blow-up caused by the change of notation. This is what is measured by succinctness.

5.2.3 Succinctness

Succinctness (Def. 5.5) actually allows to compare the size of the diagram before and after translation.

Definition 5.5 (Succinctness) *Let \mathcal{G} be a set of functions from $\mathbb{N} \rightarrow \mathbb{N}$. A language \mathcal{L}_1 is \mathcal{G} -as succinct as \mathcal{L}_2 , noted $\mathcal{L}_2 \leq \mathcal{G}(\mathcal{L}_1)$, iff there is a translation $T : \mathcal{L}_1 \rightarrow \mathcal{L}_2$ that is within \mathcal{G} : $\exists g \in \mathcal{G}, \forall n \in \mathbb{N}, \forall l_1 \in \mathcal{L}_1, |l_1| \leq n \Rightarrow |T(l_1)| \leq g(n)$. Common values for \mathcal{G} are “identically” = $\{n\}$, “thrice” = $\{3n\}$, “linearly” = $O(n)$, “cubically” = $O(n^3)$, “exponentially” = $O(2^n)$.*

By definition, wherever there is an embedding, there also exists a linear translation. In our case, a vDFD produced from a tree-shaped FD is identically as succinct as the tree, and a vDFD produced from a restricted DAG is linearly as succinct as the latter (because intermediate nodes and edges need to be added). Also, the translation from a vDFD to an RFD is linear. In all those cases, the transformation engines do not face tractability issues. However, for turning unrestricted DAG into vDFD, we need to precompute all shared cases that vDFD will treat as independent. This will cause an exponential blow-up. In practice, this means that one will be able to apply such transformations only to small diagrams.

6 Related work

Beside vDFD, a few more formally defined FD languages exist and need to be compared with FFD (and also possibly between themselves) following the same methodology:

1. Batory [1] provides a translation of FD to both grammars and propositional formulae. His goal is to use off-the-shelf Logic-Truth Maintenance Systems and SAT solvers in feature modelling tools. The semantics of grammars is a set of strings, and thus order and repetition are kept. The semantics of propositional formulae is closer to our products but [1] differs in two respects: (i) decomposable features are not eliminated, and (ii) the translation of operators by an equivalence leads to (we suspect) a counter-intuitive semantics.
2. In [8, 7], Czarnecki *et al.* define a new FD language to account for staged configuration. They introduce *feature cardinality* (the number of times a feature can be repeated in a product) in addition to the more usual (*group*) cardinality. Foremost, a new semantics is proposed where the full shape of the unordered tree is

important, including repetition and decomposable features. The semantics is defined in a 4-stage process where FD are translated in turn into an extended abstract syntax, a context-free grammar and an algebra. In [7], the authors provide an even richer syntax. The semantics of the latter is yet to be defined, but is intended to be similar to [8].

3. Benavides *et al.* [2] propose to use constraint programming to reason on feature models. They extend the FD language of [8] with extra-functional features, attributes and relations between attributes. Subsequently, they describe tool-support based on mapping those FD to Constraint Satisfaction Problems (CSP).
4. Wang *et al.* [29] propose a semantics of FD using ontologies. A semantic web environment is used to model and verify FD with OWL DL. The RACER reasoner is used to check inconsistencies during feature configuration. Their semantics slightly differ from ours, since (1) they omit justifications and (2) they did not eliminate auxiliary symbols. Missing justifications yield strongly counter-intuitive results, but keeping auxiliary symbols is harmless for consistency checking (but incorrect for other tasks) as shown in [4].

As mentioned in Sec. 1, we think that these approaches do not rank as good as ours on generality, abstraction and intuitiveness. However, a finer comparison is required in order to further justify these claims just as we have done in this paper for vDFD. This a topic for future work.

7 Conclusion

This work is a first step towards the comparative semantics of feature diagrams. We have recalled FFD, a generic formalization of feature diagrams. We have also recalled and revisited van Deursen and Klint’s definition of FD, which we called vDFD. We have then compared the two by applying the principles of comparative semantics. We have defined an abstraction function that relates van Deursen and Klint’s semantics to our own and then studied 3 properties of vDFD: expressiveness, embeddability and succinctness.

We can summarize the conclusions and practical implications of our investigations as follows:

- By being able to abstract vDFD’s semantics to ours and assuming that the abstracted information (the order of features) is irrelevant, we further validated the appropriateness of our semantic domain and semantic function for representing product lines. This gives us more assurance that FFD can be used as a reference semantics for implementing future tools for engineering the requirements of software product lines.

- We have discovered a few minor errors in the original formalization of vDFD which can cause an automatic reasoner based on it to yield erroneous results. These findings can help the developers to improve their tool and future developers to avoid the same problems.
- More fundamentally, we analyse the lack of abstraction and the presence of errors in the original formalization as a result of the bias imposed by the tool. This tends to confirm the advantages of our methodology [26, 25] where we recommend tool-independent formalization prior to the adoption or development of any supporting tool. The generic semantic framework of FFD greatly facilitates such a definition (which can boil down to filling a simple 4-entry template as shown in Table 3). From there on, the available results of comparative analyses can guide the adopters in the selection or development of their supporting tools.
- We have proved that vDFD are expressively complete with respect to the semantic domain of product lines, just as some other members of the FFD family. From this, we have concluded that the vDFD language and its supporting tools can be used to model, and reason on, product lines without the fear of being limited in expressiveness. However, this conclusion has to be mitigated by the preceding conclusions concerning the lack of abstraction and presence of errors in the original tool-driven formalization.
- By looking at translations, we have assessed the feasibility of transforming vDFD models into other kinds of feature models and conversely. Based on these results, developers can start writing model transformation scripts to enable tool interoperability. For example, if some reasoning facilities are not supported by tool A but they are supported by tool B, embeddability of A's notation into B's guarantees that a transformations from A to B is not only possible but also that B preserves the shape of the original model so that traceability between the two models is greatly facilitated. Succinctness measures the cost of the translations. Regarding translations that have vDFD as a target, we have observed that, roughly stated, the translation from tree-shaped FD are easily tractable whereas those from DAG-shaped FD are only tractable for small diagrams. Finally, vDFD have been found embeddable into RFD.

Comparative semantics should be obtained for all variants of feature diagrams, so that we know precisely what they are, what are their respective qualities: expressiveness, succinctness, embeddability, but also axiomatizability, complexity of the reasoning tasks, etc. These objective data can then serve to guide the selection of a common standard for feature diagrams. Such a standard could hopefully

accelerate their adoption by industry, increase the competition between tool providers, and help to establish efficient and safe requirements and software product line engineering methods.

8 Threats to validity

The main threats to the validity of our method to compare FD languages is that it considers only the formal aspects of these languages. Hence, we are faced with the usual advantages and drawbacks of any formalization endeavour [13]. On the one hand, a mathematically defined semantics tends (i) to eliminate ambiguity in the interpretation of languages, (ii) to facilitate safe and efficient tool automation and (iii) to allow the definition of objective criteria to evaluate and compare languages. On the other hand, a formal semantics can never guarantee by itself that it captures enough and the right information about the domain being modeled.

More general language evaluation frameworks exist that take into account a wider spectrum of qualities and criteria. For example, Krogstie [18] proposes a comprehensive language and model quality framework covering such notions as *domain appropriateness* (i.e. is the language considered adequate to convey the relevant information on the subject domain). A complete evaluation and comparison of FD languages should also take such aspects into account. These are topics for future work which need to be addressed by empirical validation when we will have developed a concrete (user-oriented) syntax and supporting tools based on FFD. Then, we will be able to evaluate whether there is a good match between the intended ("real world") and the formal semantic domains and functions of our language. Yet, it should be noted that even if we obtain a good assurance that there is such a match, we will still not be able to guarantee that our language will always be used to represent the relevant information about the requirements of the system at hand for no formal, nor informal, language can ever guarantee this by itself [14].

Another threat to the validity of our results is that all the formal definitions and reasoning in this paper have been carried out by humans without the assistance of tools (except text editing tools). Therefore, we cannot guarantee that human errors are completely absent from our comparative analysis.

Finally, it should be noted that we have only assessed the language defined by van Deursen and Klint from the description that is made of it in [27]. We have not assessed the tool itself, which might well have been improved to address the issues that we mention.

References

- [1] D. S. Batory. Feature Models, Grammars, and Propositional Formulas. In Obbink and Pohl [20], pages 7–20.
- [2] D. Benavides, P. Trinidad, and A. R. Cortés. Automated Reasoning on Feature Models. In O. Pastor and J. F. e Cunha, editors, *CAiSE*, volume 3520 of *Lecture Notes in Computer Science*, pages 491–503. Springer, 2005.
- [3] Y. Bontemps, P. Heymans, P.-Y. Schobbens, and J.-C. Trigaux. Semantics of FODA Feature Diagrams. In T. Männistö and J. Bosch, editors, *Proc. of Workshop on Software Variability Management for Product Derivation (Towards Tool Support)*, Boston, August 2004.
- [4] Y. Bontemps, P. Heymans, P.-Y. Schobbens, and J.-C. Trigaux. Generic Semantics of Feature Diagrams Variants. In *Proceedings of 8th International Conference on Feature Interactions in Telecommunications and Software Systems(ICFI)*. IOS Press, 2005.
- [5] K. D. Bosschere and J.-M. Jacquet. Comparative semantics of μ Log. In D. Etiemble and J.-C. Syre, editors, *Proceedings of PARLE '92 – Parallel Architectures and Languages Europe*, Lecture Notes in Computer Science, pages 911–926, Paris, June 15–18, 1992. Springer-Verlag.
- [6] S. Cohen, B. Tekinerdogan, and K. Czarnecki. A case study on requirement specification: Driver Monitor. In *Workshop on Techniques for Exploiting Commonality Through Variability Management at the Second International Conference on Software Product Lines (SPLC2)*, 2002.
- [7] K. Czarnecki and S. H. ad Ulrich Eisenecker. Staged Configuration Using Feature Models. *Software Process Improvement and Practice, special issue on Software Variability: Process and Management*, 10(2):143 – 169, 2005.
- [8] K. Czarnecki, S. Helsen, and U. W. Eisenecker. Formalizing cardinality-based feature models and their specialization. *Software Process: Improvement and Practice*, 10(1):7–29, 2005.
- [9] J. W. de Bakker and J. N. Kok. Comparative metric semantics for Concurrent Prolog. *Theoretical Computer Science*, 75(1–2):14–43, 25 Sept. 1990.
- [10] U. W. Eisenecker and K. Czarnecki. *Generative Programming: Methods, Tools, and Applications*. Addison-Wesley, 2000.
- [11] M. Eriksson, J. Börstler, and K. Borg. The PLUSS Approach - Domain Modeling with Features, Use Cases and Use Case Realizations. In Obbink and Pohl [20], pages 33–44.
- [12] M. Griss, J. Favaro, and M. d'Alessandro. Integrating Feature Modeling with the RSEB. In *Proceedings of the Fifth International Conference on Software Reuse*, pages 76–85, Vancouver, BC, Canada, June 1998.
- [13] D. Harel and B. Rumpe. Meaningful Modeling: What's the Semantics of "Semantics"? *IEEE Computer*, 37(10):64–72, October 2004.
- [14] M. Jackson. Some Basic Tenets of Description. *Software and Systems Journal*, 1(1):5–9, September 2002.
- [15] D. Janssens and G. Rozenberg. On the structure of node label controlled graph languages. *Inform. Sci.*, 20:191–244, 1980.
- [16] K. Kang, S. Cohen, J. Hess, W. Novak, and S. Peterson. Feature-Oriented Domain Analysis (FODA) Feasibility Study. Technical Report CMU/SEI-90-TR-21, Software Engineering Institute, Carnegie Mellon University, Nov. 1990.
- [17] K. C. Kang, S. Kim, J. Lee, and K. Kim. FORM: A Feature-Oriented Reuse Method. In *Annals of Software Engineering* 5, pages 143–168, 1998.
- [18] J. Krogstie. A Semiotic Approach to Quality in Requirements Specifications. In *Organizational Semiotics*, pages 231–249, 2001.
- [19] T. Mossakowski. Comorphism-based grothendieck logics. In K. Diks and W. Rytter, editors, *MFCS*, volume 2420 of *Lecture Notes in Computer Science*, pages 593–604. Springer, 2002.
- [20] J. H. Obbink and K. Pohl, editors. *Software Product Lines, 9th International Conference, SPLC 2005, Rennes, France, September 26-29, 2005, Proceedings*, volume 3714 of *Lecture Notes in Computer Science*. Springer, 2005.
- [21] K. Pohl, G. Bockle, and F. van der Linden. *Software Product Line Engineering: Foundations, Principles and Techniques*. Springer, july 2005.
- [22] M. Riebisch. Towards a More Precise Definition of Feature Models. *Position Paper. In: M. Riebisch, J. O. Coplien, D. Streitferdt (Eds.): Modelling Variability for Object-Oriented Product Lines*, 2003.
- [23] M. Riebisch, K. Böllert, D. Streitferdt, and I. Philippow. Extending Feature Diagrams with UML Multiplicities. In *Proceedings of the Sixth Conference on Integrated Design and Process Technology (IDPT 2002)*, Pasadena, CA, June 2002.
- [24] M. Ryan and P. Y. Schobbens. FireWorks: A Formal Transformation-Based Model-Driven Approach to Features in Product Lines. In T. Männistö and J. Bosch, editors, *Proc. WS. on Software Variability Management for Product Derivation - Towards Tool Support*, 2004.
- [25] P.-Y. Schobbens, P. Heymans, J.-C. Trigaux, and Y. Bontemps. Feature Diagrams: A Survey and A Formal Semantics (in press). In *Proceedings of the 14th IEEE International Requirements Engineering Conference (RE'06)*, in press, Minneapolis, Minnesota, USA, Sept. 2006.
- [26] P.-Y. Schobbens, P. Heymans, J.-C. Trigaux, and Y. Bontemps. Generic Semantics of Feature Diagrams (in press). *Special issue of Computer Networks on feature interactions in emerging application domains*, 2006.
- [27] A. van Deursen and P. Klint. Domain-Specific Language Design Requires Feature Descriptions. *Journal of Computing and Information Technology*, 2001.
- [28] J. van Gurp, J. Bosch, and M. Svahnberg. On the Notion of Variability in Software Product Lines. In *Proceedings of the Working IEEE/IFIP Conference on Software Architecture (WICSA'01)*, 2001.
- [29] H. Wang, L. Y. Fang, J. Sun, H. Zhang, and J. Z. Pan. A Semantic Web Approach to Feature Modeling and Verification. In *Proc. of the International Workshop on Semantic Web Enabled Software Engineering (SWESE)*, 2005.